Maximum Likelihood Combining in Cooperative Relay Networks with Different Modulations

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Abstract— This paper proposes a novel combining strategy with different modulations for cooperative relay networks. We model a Maximum Likelihood (ML) detection at the destination with higher order modulations such as quadrature amplitude modulation (QAM) which accounts individual symbol error rate (SER) to facilitate the detection. Our proposed algorithm is flexible to signals with different modulations as detection is done symbol-bysymbol basis. If different modulations are used at the source and the relays, we propose that lower modulation is used at the source. By computer simulations, significant Packet Error Rate (PER) performance can be obtained by the proposed scheme against Cooperative-Maximum Ratio Combining (C-MRC).

Keywords- relay networks; diversity combining; symbol error rate.

I. INTRODUCTION

Various relaying schemes have been proposed to explore the benefits of cooperative networks, including Decode-and-Forward (DF), Amplify-and-Forward (AF) [1]-[3] and Detect-and-Forward (DEF). DEF is simple in complexity where a relay detects the signals, modulates before forwarding to the destination. However, with no error protection, the forwarded symbols may be incorrect and, thus, the error probability must be considered at the destination for optimal detection. Recently, various combining schemes have been investigated in [4]-[9] to minimize the impact of these errors. Our previous work on Maximum Likelihood (ML) detection [4] simplifies the conventional ML-based detection which approximates the source-relay (S-R) link at sufficiently high signal-to-noiseratio (SNR). The authors in [5] developed a detector approximating the ML detection. However, this scheme cannot achieve full diversity for multi relay schemes. In [6], Cooperative-MRC (C-MRC) is introduced and yet, it suffers losses in assymetrical networks and is not easily feasible in relay networks with different modulations. In [7]-[8], the authors have presented a combining strategy with perfect channel state information (CSI) of S-R link at the destination. However, their work assumes average symbol error rate (SER) as the side information and hence, this strategy does not offer a complete ML solution. To solve this problem, we proposed an ML-based combining

strategy in [9] which exploits individual SER for the detection at the destination with quadrature phase-shift keying (QPSK) signals. With this information, we can accurately model the transition probabilities for the erroneous transmission from noisy relay channels. Nonetheless, [4]-[9] assume the modulation used by the source and the relays to be the same. In some favorable conditions, the source can use higher power and larger symbol constellations to optimize the channel resources. One finds that a conventional maximal ratio combining (MRC) cannot be used for combining signals received in different modulation formats. Selection combining (SC) was proposed for this purpose [10] and yet, this strategy is far from being optimal. In [11], Soft-bit MRC is introduced which performs well with different modulations and the application is limited to perfect relays.

In this paper, we extend our work in [9] and generalize it to quadrature amplitude modulations (QAM). Here, we present an ML combining scheme which exploits effectively perfect knowledge of all links for optimal combining at the destination and provides a solution for combining noisy relayed signals with different modulation levels. Through computer simulations we observe that the proposed scheme is not only practical to the different modulated signals but also shows a remarkable potential in achieving significant diversity gains with better packet error rate (PER) performance than that of C-MRC.

The structure of this paper is as follows: Section II is System Description and the proposed scheme, Simulation Results and Discussions are given in Section III, and in IV the paper is summarized. The derivation of the individual SER for 16QAM is presented in the Appendix.

II. SYSTEM DESCRIPTION

A. System Model

We consider a general case shown in Fig. 1, a source node (S) and a destination (D) with L relays R_l , $l \in \{1, 2, ..., L\}$ over flat Rayleigh fading channels. Assuming time division multiplexing, the source transmits its signal x_s in timeslot 1 to the destination and the relays with the average power E_s . Due to the broadcast transmission, both the destination and all L relays receive noisy symbols of x_s .



Figure 1: Block diagram of the cooperative relay system with multiple relay channels.

The received signals at the destination and at the *l*th relay can be written respectively as

$$y_{sd} = h_{sd} x_s + n_{sd}$$

$$y_{sr,l} = h_{sr,l} x_s + n_{sr,l}$$
(1)

where the subscripts indicate the node relation such that h_{sd} and $h_{sr,l}$ are independent complex-valued channel gains for the S-D link and S-R link of the *l*th relay respectively. For simplicity, all channels are Quasi-static Rayleigh fading channels i.e., $h_{sd} \sim CN(0,1)$ and $h_{sr,l} \sim \mathcal{CN}(0,1)$, where $\mathcal{CN}(\mu,\sigma^2)$ denotes a complex Gaussian random variable with mean μ and variance σ^2 . n_{sd} and n_{srl} are independent additive white Gaussian noise at the destination and the relay respectively which are modeled as $n_{sd} \sim \mathcal{CN}\left(0, \sigma_{sd}^2\right), n_{sr,l} \sim \mathcal{CN}\left(0, \sigma_{sr,l}^2\right)$ with variance equal to $N_0/2$ per dimension. We assume that the average SNR for all links are the same denoted as $\overline{\gamma} = E_s / N_0$, while the instantaneous SNR is represented as $\gamma_{sd} = |h_{sd}|^2 \overline{\gamma}$ and $\gamma_{sr,l} = |h_{sr,l}|^2 \overline{\gamma}$ respectively. The relay performs a hard-decision detection (DEF) and remodulates the detected symbol as $x_{r,l}$ with the same average power E_s for re-transmissions in timeslot 2. The symbol received at the destination is given as

$$y_{rd,l} = h_{rd,l} x_{r,l} + n_{rd,l}$$
 (2)

where $h_{rd,l} \sim CN(0,1)$ and $n_{rd,l} \sim CN(0,\sigma_{rd,l}^2)$ with variance equal to $N_0/2$ per dimension. The instantaneous SNR is $\gamma_{rd,l} = |h_{rd,l}|^2 \overline{\gamma}$. At the destination, the received signals from the source and the relay node are combined in order to recover the original source data.

B. Proposed ML-based Combining Strategy

Here, we generalize our work in [9] to *M*-QAM with different modulations at the source and the relay nodes. First, let us take a closer look at C-MRC. In [6], C-MRC output at the destination is

$$y_{cmrc} = h_{sd}^* y_{sd} + \sum_{l=1}^{L} \frac{\gamma_{\min,l}}{\gamma_{rd,l}} h_{rd,l}^* y_{rd,l}$$
(3)

where $\gamma_{\min,l} = \min(\gamma_{sr,l}, \gamma_{rd,l})$, $\gamma_{sr,l}$ and $\gamma_{rd,l}$ are instantaneous SNR of the S-R and R-D channels for the *l*th relay node respectively. $\gamma_{\min,l}$ is tight approximation of the equivalent SNR of the S-R-D link at high SNR [6]. The usual intuitive meaning associated with (3) is that when $\gamma_{sr,l}$ is high, the detector places full confidence to the relayed signals. In case of low $\gamma_{sr,l}$, the confidence is weighted according to the ratio of both hops that is S-R-D link. In fact, like MRC, (3) implies that C-MRC cannot be easily used for signals with different modulations.

The algorithm in [9] optimally combines the noisy signals received at the destination node, y_{sd} and $y_{rd,l}$ by considering the effect of detection errors at the output of the *l*th relay. However, the focus is only on the combining method with the same modulation, QPSK at both the source and the relays. From [9], the corresponding joint ML decision criterion finds \hat{x}_s , an estimate of x_s and is defined as

$$\hat{x}_{s} = \arg\max_{x_{s}, x_{r} \in \{\chi_{s} = \chi_{r,l}\}} p_{sd}(y_{sd} \mid x_{s}) \times \prod_{l=1}^{L} \left\{ P(x_{r,l} = x_{s}) p_{rd}(y_{rd,l} \mid x_{r,l} = x_{s}) + \sum_{x_{r,l} \in \chi_{r,l}} P(x_{r,l} \neq x_{s}) p_{rd}(y_{rd,l} \mid x_{r,l}) \right\}$$
(4)

where χ_s and $\chi_{r,l}$ denotes the finite set of the constellation at the source and the *l*th relay respectively; we use capital *P* as the probability ; $p_{sd}(y_{sd} | x_s)$ is the PDF of the source signal y_{sd} conditioned upon the transmitted signal x_s and $p_{rd}(y_{rd,l} | x_{r,l} = x_s)$ is the PDF of the relayed signal $y_{rd,l}$ conditioned on the equality of both transmitted symbols $(x_{r,l} = x_s)$. The bracketed term in (4) has to consider the error probability of the received signal $y_{sr,l}$ at the *l*th relay accounting the individual SER of each signal point.

In QPSK, the transmit symbol x_s which is labeled by two bits, (b_1, b_2) takes from the constellation set $\chi_s \in \{s_1, s_2, s_3, s_4\}$. Assuming the source and the relays use the same QPSK modulation i.e., $\chi_s = \chi_{r,l}$, the detection at the destination is performed jointly by the ML criterion and we can expand (4) as

$$\hat{x}_{s} = \arg \max_{x_{s}, x_{r} \in \{\chi_{s} = \chi_{r,l}\}} p_{sd}(y_{sd} \mid x_{s}) \times \left\{ \prod_{l=1}^{L} \begin{cases} [1 - (\varepsilon_{l} + 2\varepsilon_{2})] p_{rd}(y_{rd,l} \mid x_{r,l} = x_{s}) \\ + \varepsilon_{1} p_{rd}(y_{rd,l} \mid x_{r,l} = x_{s} \varepsilon^{j\pi}) \\ + \varepsilon_{2} p_{rd}(y_{rd,l} \mid x_{r,l} = x_{s} \varepsilon^{j\frac{\pi}{2}}) \\ + \varepsilon_{3} p_{rd}(y_{rd,l} \mid x_{r,l} = x_{s} \varepsilon^{-j\frac{\pi}{2}}) \end{cases} \right\}$$
(5)

where ε_1 , ε_2 and ε_3 denote the symbol error probabilities from $s_1 \rightarrow s_3$, $s_1 \rightarrow s_2$ and $s_1 \rightarrow s_4$ respectively; ε_1 and $\varepsilon_2 = \varepsilon_3$ are analytically expressed as the Gaussian Qfunction where $Q(x) = (1/\sqrt{2\pi}) \int_x^{\infty} \exp(-t^2/2) dt$. In the bracketed term of (5), we include the multiplicative error term in exponential function, $\varepsilon^{j\phi}$ with the following equality $x_{r,l} = x_s \varepsilon^{j\phi}$ where $\phi \in \{0, \pi, \frac{\pi}{2}, -\frac{\pi}{2}\}$ denotes the phase changes that depends on the symbols transmitted from the relay. This means that in (5), the relay does not operate error-free. In this paper, we employ closed-form expressions for the probability of error for each constellation symbol for QPSK as

$$\varepsilon_{1} = P(x_{r,l} = x_{s}\varepsilon^{j\pi})$$

$$= \frac{1}{4} \operatorname{erfc}^{2} \left[\sqrt{\frac{E_{s}}{2N_{0}}} \right] = \left\{ Q\left(\sqrt{\frac{E_{s}}{N_{0}}}\right) \right\}^{2}$$

$$\varepsilon_{2} = P(x_{r,l} = x_{s}\varepsilon^{j\frac{\pi}{2}}) = \varepsilon_{3} = P(x_{r,l} = x_{s}\varepsilon^{-j\frac{\pi}{2}})$$

$$= Q\left(\sqrt{\frac{E_{s}}{N_{0}}}\right) - \left\{ Q\left(\sqrt{\frac{E_{s}}{N_{0}}}\right) \right\}^{2}$$

$$(6)$$

$$(7)$$

where erfc is the complementary error function. When y_{sd} and $y_{rd,l}$ are received at the destination, by inserting s_1, s_2, s_3 or s_4 to x_s and examining how large the argument value in (5), we can determine the transmit signal point x_s from the finite set χ_s in QPSK constellation. The PDF expression in (5) can be represented by

$$p_{rd}(y_{rd,l} | x_{r,l} = x_s \varepsilon^{j\varphi}) = \frac{1}{\sqrt{2\pi\sigma_{rd,l}^2}} \exp\left\{-\frac{|y_{rd,l} - h_{rd,l} x_s \varepsilon^{j\varphi}|^2}{2\sigma_{rd,l}^2}\right\}$$
(8)

where $\phi \in \{0, \pi, \pi/2, -(\pi/2)\}\)$. The analytical results presented thus far in previous works have been derived from studies which examined the SER problem assuming that the SER of each QPSK symbol is equally likely (average SER). Thus, these results cannot be treated as offering a complete ML solution. Note that another advantage in the proposed ML over C-MRC is its flexibility of combining different modulated signals from different nodes since each link can be treated independently (symbol-wise detection).

Next, we generalize (5) to *M*-QAM. From here, we observe that there are $\chi_{r,l}$ –1 ways of making an incorrect decision and their impacts on detection at the destination are not necessarily the same. Thus, we can easily show the decision criterion for general *M*-QAM as

$$\hat{x}_{s} = \arg \max_{x_{s}, x_{r} \in \{\chi_{s} \neq \chi_{r,l}\}} p_{sd} (y_{sd} \mid x_{s}) \times \left\{ \prod_{l=1}^{L} \left\{ (1 - \sum_{\kappa=1:\chi_{r,l}-1} \varepsilon_{\kappa}) p_{rd} (y_{rd,l} \mid x_{r,l} = x_{s}) + \sum_{\kappa=1:\chi_{r,l}-1} \varepsilon_{\kappa} p_{rd} (y_{rd,l} \mid x_{r,l}^{\kappa}) \right\}$$
(9)

where ε_{κ} , $\kappa = \{1, 2, ..., \chi_{r,l} - 1\}$ is the SER for each symbol in *M*-QAM based on the modulation size in each relayed path and are expressed in Q-function as well. For example in the Appendix, we illustrate the derivations of ε_{κ} for some 16QAM symbols.

C. Analysis of Combining Schemes

Now, we analyze the channel capacity of C-MRC and the proposed ML schemes. Here, we assume one relay node for simplicity. Let us denote the channel capacities of S-R, R-D and S-D links by $C_{sr}(\gamma_{sr}) = \log_2(1+\gamma_{sr})$, $C_{rd}(\gamma_{rd}) = \log_2(1+\gamma_{rd})$ and $C_{sd}(\gamma_{sd}) = \log_2(1+\gamma_{sd})$ respectively, and the joint capacity of the combined signals at the destination during the cooperative phase by C_{tol} . The channel capacity unit is bit per channel use. The total capacities C_{tol} for C-MRC and the proposed scheme are as follows [12]

$$C_{tol}^{C-MRC} = C(\gamma_{sd} + \gamma_{rd})$$

$$C_{tol}^{ML} = C(\gamma_{sd}) + C(\gamma_{rd})$$
(10)

Assuming the instantaneous SNR for each link $\gamma_{sr}, \gamma_{rd}, \gamma_{sd} \ge 0$ then we have

$$C(\gamma_{sr} + \gamma_{rd}) + C(\gamma_{sd}) \ge C(\gamma_{sd} + \gamma_{rd}) + C(\gamma_{sr})$$
(11)

if and only if $(\gamma_{sd} \ge \gamma_{sr} \lor \gamma_{rd} = 0)$. (10) and (11) show that the variations in the relayed link reduces the total channel capacity. Particularly, the degradation in performance of C-MRC can be worse than that of the ML i.e., $C_{tol}^{ML} \ge C_{tol}^{C-MRC}$. We also prove this claim by computer simulations in what follows.

D. Complexity Comparison

The computational complexity of the receiver at the destination depends on the detection algorithms, the hardware architectures, and other factors. In this paper, we evaluate the computational complexity for our proposed scheme, C-MRC and SC based on the number of complex multiplications and additions. For convenience, we consider the required computations for the functions of equalization, detection and signal combining at the destination in a relay node scheme (L = 1) only. Here we assume QPSK modulation is used at the source and relay node. We define that each multiplication from two complex numbers takes four complex multiplications and two additions. If Euclidean distance metric calculation is employed, we need 46 complex multiplications and 16 additions to detect a symbol at the receiver. Thus, this becomes the baseline computational complexity for SC strategy. Due to space limitation, other derivations are omitted for brevity. Table 1 compares the number of required complex multiplications and additions for each scheme per symbol.

Table 1: The number of complex multiplications and additions at each scheme.

Complexity	SC	C-MRC	Proposed
Multiplication	46	50	230
Addition	16	28	80

Table 1 shows that the computational complexity increases with the order from SC<C-MRC<ML-based Combining (proposed). SC turns out to be the lowest but with a significant reduction in the error rate performance as shown in the following section. SC only uses one signal for detection at the receiver and hence, the computation is less. This outcome for our proposed scheme is expected since the additional complexity in the scheme is coupled with a significant error rate improvement compared against the conventional SC and C-MRC in various simulation setups as shown in the manuscript. The complexity of the proposed scheme is highest because the destination has to consider individual SER of making wrong decisions at relay nodes in the detection. Thus, the complexity of the proposed scheme increases as the modulation increases. However, to assist the detection at the destination, our proposed scheme only requires the average receive SNR of S-R link to compute individual SER of the modulation as shown in (6) and (7). Therefore, our proposed scheme still inherits an interesting trade-off between the error rate performance and the system complexity. Although C-MRC is simpler in the computational complexity, its biggest challenge is to have accurate instantaneous channel knowledge at the receiver. In practice, one needs accurate channel estimation and a high signaling overhead in C-MRC scheme to feedback the channel knowledge to the destination. In fact, there is no practical C-MRC approach ever proposed yet for combining different modulated signals.

III. SIMULATION RESULTS AND DISCUSSIONS

We simulate PER against average SNR in decibel (dB). For convenience, we restrict our work to QPSK and 16QAM modulations only. To reduce the computational complexity in the proposed ML for 16QAM, we adopt the max-log approximation. We assume the source and all relay nodes transmit with the same average power E_s resulting in the average SNR, $\overline{\gamma} = E_s / N_0$ (symmetrical network). For C-MRC, we also consider the destination has a perfect knowledge of S-R link (i.e., instantaneous SNR) and perfect channel estimation is assumed.

Fig. 2 shows the PER of the proposed scheme for 16QAM modulation at both the source and the relays, $\chi_s = \chi_{r,l}$ against the baseline for multiple relay nodes i.e., L = 1, 2 and 3. As expected, the proposed schemes outperform C-MRC (3) in all cases with 0.5dB, 1dB and 1.5dB gap at PER = 10^{-3} for 1, 2 and 3 relay cases respectively. We notice that all of the cooperative schemes achieve full diversity order as viewed from the slopes of the curves i.e., $10^{-(L+1)}/10$ (dB) (diversity order $\overline{\gamma}$ of L+1). This result shows that the proposed algorithm has better accuracy of detection due to the sufficient statistics of the received signals y_{sd} and $y_{rd,l}$. For this reason, the conditional probability $p_{rd}(y_{rd,l} | x_{r,l})$ can be computed using the observations $p_{rd}(y_{rd,l} | x_{r,l} \neq x_s)$.

Next, in Fig. 3 we simulate the proposed scheme and C-MRC with 16QAM in both nodes under different R-D link quality. From here onwards we only simulate for case (L = 1). Thus, for convenience we remove the subscript l in the notation. We vary the average SNR for R-D link $\overline{\gamma}_{rd}$, and we keep the average SNR for S-D link and S-R link the same, $\overline{\gamma}_{sd} = \overline{\gamma}_{sr} = \overline{\gamma}$. We simulate the schemes at three different scenarios of R-D link quality: $\overline{\gamma}$ +15dB (+15dB), $\overline{\gamma}$ -15dB (-15dB) and $\overline{\gamma}_{rd} = \overline{\gamma}$ (equal). In Fig. 3, we find that the proposed scheme can outperform C-MRC when R-D link has sufficiently high SNR quality (+15dB) with marginal 1dB gap at PER= 10^{-3} and 2.5dB gap at $PER=10^{-2}$ for low SNR quality (-15dB). When R-D link has higher SNR compared to S-D link, the combined signal at the destination is dominated by the errors from the relayed link whose error is due to the detection error at the relay. Given that the relay has made a decision error and hence, the source and the relay send contradicting information to the destination. As a result, when the R-D has very low SNR, the PER performance is degraded further compared to the case of equal SNR. In C-MRC, one can also refer to (3) that C-MRC effectiveness is largely conditioned on the link quality of R-D link over S-D link (direct path). In particular, the proposed scheme improves achievable PER performance which becomes an added advantage compared to C-MRC. This result also confirms the channel capacity analysis in II-C.



Figure 2: PER comparison with 16QAM at the source and the relay nodes between the proposed ML scheme and C-MRC (dashed lines) using DEF protocols for L = 1, 2 and 3 relays.



Figure 3: PER comparison with 16QAM at the source and the relay node between the proposed ML scheme (solid lines) and C-MRC (dashed lines) when the average SNR of R-D link, $\overline{\gamma}_{rd}$ varies at L = 1 relay case.



Figure 4: PER comparison proposed ML scheme and SC with 16QAM at the source and different modulation at the relay using DEF protocols for L = 1 relay.

In Fig. 4, we simulate the proposed scheme with different modulations, QPSK and 16QAM at the relay node (L = 1). For comparison, we use selection combining

(SC) with the same channel setup. We also simulate a scheme when no relay is used with BPSK modulation. A simulated lower bound with one perfect relay (i.e., errorless relay detection) is included in this simulation. The results in Fig. 4 clearly show that the proposed scheme outperforms SC scheme with great margins. In both combining techniques, as expected, we observe that there is a slight improvement in PER if lower modulation i.e., QPSK is used at the relay which is about 1dB gain at PER= 10^{-2} . This result is expected due to the fact that lower modulation is less vulnerable to errors.





However, the result in Fig. 4 does not consider the same total transmission rate at the destination. In Fig. 5, in transmission а fixed rate scheme i.e., $\eta = ((\log_2 \chi_s)^{-1} + (\log_2 \chi_r)^{-1})^{-1}$ which means that η is the same for the cases in comparison, we simulate when the source and the relay use different modulation assignments, $\chi_s \neq \chi_r$. For simplicity, the scheme uses 2 sets of modulation combinations from QPSK and 16QAM. In the proposed scheme, since different modulations carry different number of bits per symbol, we propose to do bitby-bit detection if mapping conversion is required at the relay node. To extract the bits from the symbols, symbol log-likelihood ratio (LLR) can be used [4]. Thus, regardless of the modulation constellations used at the relay, we can easily convert the mapping from QPSK to higher constellations or vice versa. For case 1 when S uses OPSK, the relay employs 16OAM (S=OPSK, R=16OAM). In case 2, the source uses 16OAM and the relay uses QPSK (S=16QAM, R=QPSK) which is identical to the curves in Fig. 4. From Fig. 5, the result clearly shows that the proposed scheme performs better when lower modulation is used at the source which is about 3dB improvement by the proposed scheme at PER $=10^{-3}$. The proposed scheme easily achieves the full diversity gain of 2 for both cases. The same trend occurs in SC scheme with around 3dB improvement at PER=10⁻² but with lower diversity gain due to the error propagation from the relay. The 1dB loss in the simulation result is the direct outcome of the error propagation of the noisy channels. In brief, we suggest that assigning lower modulation at the source is a better strategy to bring more performance improvement in relay networks.

IV. CONCLUSIONS

In this paper, an extension of ML-based combining strategy for cooperative relay scheme to different modulations is proposed. With the potential errors at the relays, we can accurately model the transition probabilities for the erroneous transmission from noisy relays. Our work also investigates the PER performance when the source and the relays have different modulations. We found that it is better to use lower modulation at the source, thus reducing possible error propagation from the relays. By computer simulation, the proposed ML scheme is superior to C-MRC in PER performance compared to C-MRC regardless of the modulation schemes.

APPENDIX

Derivation of Individual Symbol Error Rate (SER) of 16QAM Signals in Gray Mapping



Figure 6: 16QAM symbols and symbol error probabilities.

Here, we present the derivations of individual SER of 16QAM symbols in AWGN channels which become the side information to our proposed scheme (9). Fig. 6 depicts the signal points for 16QAM with its decision boundaries as the dashed lines when Gray mapping is used. The constellation points of 16QAM are normalized with the factor $a = 1/\sqrt{10}$ to ensure that the average energy over all symbols is unity. Let us denote *I* and *Q* as the in-phase and quadrature components respectively. Since each complex symbol of 16QAM corresponds to four binary bits, $(b_1b_2b_3b_4)$ as presented in Fig. 6 we label the respective symbols accordingly.

Similar in QPSK [9], first we consider, for instance, the symbol (0101) is transmitted from the source assuming the perfect CSI is available at the receiver side. If the receiver wrongly detects the symbol as (0001), the SER for this particular symbol is calculated as

$$\varepsilon_{2} = \int_{0}^{\frac{2}{\sqrt{10}}} \int_{\frac{2}{\sqrt{10}}}^{\infty} \frac{1}{\sqrt{2\pi\sigma^{2}}} \exp\left(-\frac{\left(x - \frac{1}{\sqrt{10}}\right)^{2} + \left(y - \frac{1}{\sqrt{10}}\right)^{2}}{2\sigma^{2}}\right) dxdy$$

$$= \frac{1}{2} erf\left[\frac{1}{2\sqrt{5\sigma}}\right] erfc\left[\frac{1}{2\sqrt{5\sigma}}\right]$$

$$= Q\left(\sqrt{\frac{E_{s}}{5N_{0}}}\right) - 2\left\{Q\left(\sqrt{\frac{E_{s}}{5N_{0}}}\right)\right\}^{2}$$
(12)

where erf is the error function. For other symbols like (0100) and (0001), identical SER can be observed due to symmetry. Likewise, the calculation for ε_3 for symbol (0000) which is located on the right top corner of the

quadrant, can be found from the following integrations as

$$\varepsilon_{3} = \int_{\frac{2}{\sqrt{10}}}^{\infty} \int_{\frac{2}{\sqrt{10}}}^{\infty} \frac{1}{\sqrt{2\pi\sigma^{2}}} \exp\left(-\frac{(x-\frac{1}{\sqrt{10}})^{2} + (y-\frac{1}{\sqrt{10}})^{2}}{2\sigma^{2}}\right) dxdy$$
$$= \frac{1}{4} erfc^{2} [\frac{1}{2\sqrt{5\sigma}}] = \left\{ Q\left(\sqrt{\frac{E_{s}}{5N_{0}}}\right) \right\}^{2}$$

(13)

Since some symbols like symbol (1101) and (0111) as shown in Fig. 5 are identical i.e., ε_1 , computation of these SERs can be reduced. Finally, other SERs can be found straightforward in a similar fashion and they are not shown here for brevity.

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